

# Embedding the Erdős-Rényi Hypergraph into the Random Regular Hypergraph and Hamiltonicity

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## Abstract

We establish an inclusion relation between two uniform models of random  $k$ -graphs (for constant  $k \geq 2$ ) on  $n$  labeled vertices:  $\mathbb{G}^{(k)}(n, m)$ , the random  $k$ -graph with  $m$  edges, and  $\mathbb{R}^{(k)}(n, d)$ , the random  $d$ -regular  $k$ -graph. We show that if  $n \log n \ll m \ll n^k$  we can choose  $d = d(n) \sim km/n$  and couple  $\mathbb{G}^{(k)}(n, m)$  and  $\mathbb{R}^{(k)}(n, d)$  so that the latter contains the former with probability tending to one as  $n \rightarrow \infty$ . This extends some previous results of Kim and Vu about “sandwiching random graphs”. In view of known threshold theorems on the existence of different types of Hamilton cycles in  $\mathbb{G}^{(k)}(n, m)$ , our result allows us to find conditions under which  $\mathbb{R}^{(k)}(n, d)$  is Hamiltonian. In particular, for  $k \geq 3$  we conclude that if  $n^{k-2} \ll d \ll n^{k-1}$ , then a.a.s.  $\mathbb{R}^{(k)}(n, d)$  contains a tight Hamilton cycle.

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# 1 Introduction

## 1.1 Background

A  $k$ -uniform hypergraph (or  $k$ -graph for short) on a vertex set  $V = [n] = \{1, \dots, n\}$  is an ordered pair  $G = (V, E)$  where  $E$  is a family of  $k$ -element subsets of  $V$ . The *degree* of a vertex  $v$  in  $G$  is defined as

$$\deg_G(v) := |\{e \in E : v \in e\}|.$$

A  $k$ -graph is  $d$ -regular if the degree of every vertex is  $d$ . Let  $\mathcal{R}^{(k)}(n, d)$  be the family of all  $d$ -regular  $k$ -graphs on  $V$ . Throughout, we tacitly assume that  $k$  divides  $nd$ . By  $\mathbb{R}^{(k)}(n, d)$  we denote the  $d$ -regular random  $k$ -graph, which is chosen uniformly at random from  $\mathcal{R}^{(k)}(n, d)$ .

Let us recall two more standard models of random  $k$ -graphs on  $n$  vertices. For  $p \in [0, 1]$ , the *binomial random  $k$ -graph*  $\mathbb{G}^{(k)}(n, p)$  is obtained by including each of the  $\binom{n}{k}$  possible edges with probability  $p$ , independently of others. Further, for an integer  $m \in [0, \binom{n}{k}]$ , the *uniform random  $k$ -graph*  $\mathbb{G}^{(k)}(n, m)$  is chosen uniformly at random among all  $\binom{n}{k}$   $k$ -graphs on  $V$  with precisely  $m$  edges.

We study the behavior of these random  $k$ -graphs as  $n \rightarrow \infty$ . Parameters  $d, m, p$  are treated as functions of  $n$  and typically tend to infinity in case of  $d, m$ , or zero, in case of  $p$ . Given a sequence of events  $(\mathcal{A}_n)$ , we say that  $\mathcal{A}_n$  holds *asymptotically almost surely* (a.a.s.) if  $\mathbb{P}(\mathcal{A}_n) \rightarrow 1$ , as  $n \rightarrow \infty$ .

In 2004, Kim and Vu [11] proved that if  $d = o(n^{1/3}/\log^2 n)$  then there exists a coupling (that is, a joint distribution) of the random graphs  $\mathbb{G}^{(2)}(n, p)$  and  $\mathbb{R}^{(2)}(n, d)$  with  $p = \frac{d}{n} (1 - O((\log n/d)^{1/3}))$  such that

$$\mathbb{G}^{(2)}(n, p) \subset \mathbb{R}^{(2)}(n, d) \quad \text{a.a.s.} \tag{1}$$

They pointed out several consequences of this result, emphasizing the ease with which one can carry over known properties of  $\mathbb{G}^{(2)}(n, p)$  to the harder to study regular model  $\mathbb{R}^{(2)}(n, d)$ . Kim and Vu conjectured that such a coupling is possible for all  $d \gg \log n$  (they also conjectured a reverse embedding which is not of our interest here). In [7] we considered a (slightly weaker) extension of Kim and Vu's result to  $k$ -graphs,  $k \geq 3$ , and proved that

$$\mathbb{G}^{(k)}(n, m) \subset \mathbb{R}^{(k)}(n, d) \quad \text{a.a.s.} \tag{2}$$

whenever  $C \log n \leq d \ll n^{1/2}$  and  $m \sim cnd$  for some absolute large constant  $C$  and a sufficiently small constant  $c = c(k) > 0$ . Although (2) is stated for the uniform  $k$ -graph  $\mathbb{G}^{(k)}(n, m)$ , it is easy to see that one can replace  $\mathbb{G}^{(k)}(n, m)$  by  $\mathbb{G}^{(k)}(n, p)$  with  $p = m/\binom{n}{k}$  (see Section 5).

## 1.2 The Main Result

In this paper we extend (2) to larger degrees, assuming only  $d \leq cn^{k-1}$  for some constant  $c = c(k)$ . Moreover, our result implies that, provided  $\log n \ll d \ll n^{k-1}$ , we can take  $m \sim nd/k$ , that is, the embedded  $k$ -graph contains almost all edges of the regular  $k$ -graph rather than just a positive fraction, as in [7]. The new result is also valid for  $k = 2$  (for the proof of this case alone, see also [10, Section 10.3]), and thus extends (1).

**Theorem 1.** *For each  $k \geq 2$  there is a positive constant  $C$  such that if for some real  $\gamma = \gamma(n)$  and integer  $d = d(n)$ ,*

$$C \left( (d/n^{k-1} + (\log n)/d)^{1/3} + 1/n \right) \leq \gamma < 1, \quad (3)$$

*and  $m = (1 - \gamma)nd/k$  is an integer, then there is a joint distribution of  $\mathbb{G}^{(k)}(n, m)$  and  $\mathbb{R}^{(k)}(n, d)$  with*

$$\lim_{n \rightarrow \infty} \mathbb{P} \left( \mathbb{G}^{(k)}(n, m) \subset \mathbb{R}^{(k)}(n, d) \right) = 1.$$

**Remark.** *In the assumption (3) of Theorem 1 the term  $1/n$  can be omitted when  $k \leq 7$ . Indeed, the inequality of arithmetic and geometric means implies that*

$$(d/n^{k-1} + (\log n)/d)^{1/3} \geq (2/n^{(k-1)/2})^{1/3} \geq \sqrt[3]{2}/n.$$

For a given  $k \geq 2$ , a  $k$ -graph property is a family of  $k$ -graphs closed under isomorphisms. A  $k$ -graph property  $\mathcal{P}$  is called *monotone increasing* if it is preserved by adding edges (but not necessarily by adding vertices, as the example of, say, perfect matching shows).

**Corollary 2.** *Let  $\mathcal{P}$  be a monotone increasing property of  $k$ -graphs and  $\log n \ll d \ll n^{k-1}$ . If for some  $m \leq (1 - \gamma)nd/k$ , where  $\gamma$  satisfies (3),  $\mathbb{G}^{(k)}(n, m) \in \mathcal{P}$  a.a.s., then  $\mathbb{R}^{(k)}(n, d) \in \mathcal{P}$  a.a.s.*

## 1.3 Hamilton Cycles in Hypergraphs

To show a more specific application of Theorem 1 we consider Hamilton cycles in random regular hypergraphs.

For integers  $1 \leq \ell < k$ , define an  $\ell$ -overlapping cycle (or  $\ell$ -cycle, for short) as a  $k$ -graph in which, for some cyclic ordering of its vertices, every edge consists of  $k$  consecutive vertices, and every two consecutive edges (in the natural ordering of the edges induced by the ordering of the vertices) share exactly  $\ell$  vertices. (For  $\ell > k/2$  it implies, of course, that some nonconsecutive edges intersect as well.) A 1-cycle is called *loose* and a  $(k - 1)$ -cycle is called *tight*. A spanning  $\ell$ -cycle in a  $k$ -graph  $H$  is called an  $\ell$ -Hamilton cycle. Observe that a necessary condition for the existence of

an  $\ell$ -Hamilton cycle is that  $n$  is divisible by  $k - \ell$ . We will assume this divisibility condition whenever relevant.

Let us recall the results on Hamiltonicity of random regular graphs, that is, the case  $k = 2$ . Asymptotically almost sure Hamiltonicity of  $\mathbb{R}^{(2)}(n, d)$  was proved by Robinson and Wormald [14] for fixed  $d \geq 3$ , by Krivelevich, Sudakov, Vu and Wormald [12] for  $d \geq n^{1/2} \log n$ , and by Cooper, Frieze and Reed [3] for  $C \leq d \leq n/C$  and some large constant  $C$ .

Much less is known for random hypergraphs. Even for the standard models, the thresholds were found only recently. First, results on loose Hamiltonicity of  $\mathbb{G}^{(k)}(n, p)$  were obtained by Frieze [8] (for  $k = 3$ ), Dudek and Frieze [4] (for  $k \geq 4$  and  $2(k-1)|n$ ), and by Dudek, Frieze, Loh and Speiss [6] (for  $k \geq 3$  and  $(k-1)|n$ ). As usual, the asymptotic equivalence of the models  $\mathbb{G}^{(k)}(n, p)$  and  $\mathbb{G}^{(k)}(n, m)$  (see, e.g., Corollary 1.16 in [9]) allows us to reformulate the aforementioned results for the random  $k$ -graph  $\mathbb{G}^{(k)}(n, m)$ .

**Theorem 3** ([8, 4, 6]). *There is a constant  $C > 0$  such that if  $m \geq Cn \log n$ , then a.a.s.  $\mathbb{G}^{(3)}(n, m)$  contains a loose Hamilton cycle. Furthermore, for every  $k \geq 4$  if  $m \gg n \log n$ , then a.a.s.  $\mathbb{G}^{(k)}(n, m)$  contains a loose Hamilton cycle.*

From Theorem 3 and the older embedding result (2), in [7] we concluded that there is a constant  $C > 0$  such that if  $C \log n \leq d \ll n^{1/2}$ , then a.a.s.  $\mathbb{G}^{(3)}(n, d)$  contains a loose Hamilton cycle. Furthermore, for every  $k \geq 4$  if  $\log n \ll d \ll n^{1/2}$ , then a.a.s.  $\mathbb{R}^{(k)}(n, d)$  contains a loose Hamilton cycle.

Thresholds for  $\ell$ -Hamiltonicity of  $\mathbb{G}^{(k)}(n, m)$  in the remaining cases, that is, for  $\ell \geq 2$ , were recently determined by Dudek and Frieze [5] (see also Allen, Böttcher, Kohayakawa, and Person [1]).

**Theorem 4** ([5]).

- (i) *If  $k > \ell = 2$  and  $m \gg n^2$ , then a.a.s.  $\mathbb{G}^{(k)}(n, m)$  is 2-Hamiltonian.*
- (ii) *For all integers  $k > \ell \geq 3$ , there exists a constant  $C$  such that if  $m \geq Cn^\ell$  then a.a.s.  $\mathbb{G}^{(k)}(n, m)$  is  $\ell$ -Hamiltonian.*

In view of Corollary 2, Theorems 3 and 4 immediately imply the following result that was already anticipated by the authors in [7].

**Theorem 5.**

- (i) *There is a constant  $C > 0$  such that if  $C \log n \leq d \leq n^{k-1}/C$ , then a.a.s.  $\mathbb{R}^{(3)}(n, d)$  contains a loose Hamilton cycle. Furthermore, for every  $k \geq 4$  there is a constant  $C > 0$  such that if  $\log n \ll d \leq n^{k-1}/C$ , then a.a.s.  $\mathbb{R}^{(k)}(n, d)$  contains a loose Hamilton cycle.*
- (ii) *For all integers  $k > \ell = 2$  there is a constant  $C$  such that if  $n \ll d \leq n^{k-1}/C$  then a.a.s.  $\mathbb{R}^{(k)}(n, d)$  contains a 2-Hamilton cycle.*

(iii) For all integers  $k > \ell \geq 3$  there is a constant  $C$  such that if  $Cn^{\ell-1} \leq d \leq n^{k-1}/C$  then a.a.s.  $\mathbb{R}^{(k)}(n, d)$  contains an  $\ell$ -Hamilton cycle.

## 1.4 Structure of the Paper

In the following section we define a  $k$ -graph process  $(\mathbb{R}^{(k)}(t))_t$  which reveals edges of the random  $d$ -regular  $k$ -graph one at a time. Then we state a crucial Lemma 6, which says, loosely speaking, that unless we are very close to the end of the process, the conditional distribution of the  $(t + 1)$ -th edge is approximately uniform over the complement of  $\mathbb{R}(t)$ . Based on Lemma 6, we show that a.a.s.  $\mathbb{G}^{(k)}(n, m)$  can be embedded in  $\mathbb{R}^{(k)}(n, d)$ , by refining a coupling similar to the one we used in [7].

In Section 3 we prove auxiliary results needed in the proof of Lemma 6. They mainly reflect the phenomenon that a typical trajectory of the  $d$ -regular process  $(\mathbb{R}(t))_t$  has concentrated local parameters. In particular, concentration of vertex degrees is deduced from a Chernoff-type inequality (the only “external” result used in the paper), while (one-sided) concentration of common degrees of sets of vertices is obtained by an application of the switching technique (a similar application appeared in [12]).

In Section 4 we prove Lemma 6. First we rephrase it as an enumerative problem (counting the number of  $d$ -regular extensions of a given  $k$ -graph). We avoid usual difficulties of asymptotic enumeration by dealing with *relative* enumeration, that is, estimating the ratio of the numbers of extensions of two  $k$ -graphs which differ just in two edges. For this we define two random multi- $k$ -graphs (via the configuration model) and couple them using yet another switching.

## 2 Proof of Theorem 1

We often drop the superscript in notations like  $\mathbb{G}^{(k)}$  and  $\mathbb{R}^{(k)}$  whenever  $k$  is clear from the context.

Let  $K_n$  denote the complete  $k$ -graph on vertex set  $[n]$ . Recall the standard  $k$ -graph process  $\mathbb{G}(t), t = 0, \dots, \binom{n}{k}$  which starts with the empty  $k$ -graph  $\mathbb{G}(0) = ([n], \emptyset)$  and at each time step  $t \geq 1$  adds an edge  $\varepsilon_t$  drawn from  $K_n \setminus \mathbb{G}(t-1)$  uniformly at random. We treat  $\mathbb{G}(t)$  as an *ordered  $k$ -graph* (that is, with an ordering of edges) and write

$$\mathbb{G}(t) = (\varepsilon_1, \dots, \varepsilon_t), \quad t = 0, \dots, \binom{n}{k}.$$

Of course, the random uniform  $k$ -graph  $\mathbb{G}(n, m)$  can be obtained from  $\mathbb{G}(M)$  by ignoring the ordering of the edges.

Our approach is to represent  $\mathbb{R}(n, d)$  as an outcome of another  $k$ -graph process which, to some extent, behaves similarly to  $(\mathbb{G}(t))_t$ . For this, generate a random

$d$ -regular  $k$ -graph  $\mathbb{R}(n, d)$  and choose an ordering  $(\eta_1, \dots, \eta_M)$  of its

$$M := \frac{nd}{k}$$

edges uniformly at random. Revealing the edges of  $\mathbb{R}(n, d)$  in that order one by one, we obtain a *regular  $k$ -graph* process

$$\mathbb{R}(t) = (\eta_1, \dots, \eta_t), \quad t = 0, \dots, M.$$

For every ordered  $k$ -graph  $G$  with  $t$  edges and every edge  $e \in K_n \setminus G$  we clearly have

$$\mathbb{P}(\varepsilon_{t+1} = e \mid \mathbb{G}(t) = G) = \frac{1}{\binom{n}{k} - t}.$$

This is not true for  $\mathbb{R}(t)$ , except for the very first step  $t = 0$ . However, it turns out that for the most of the time conditional distribution of the next edge in the process  $\mathbb{R}(t)$  is approximately uniform, which is made precise by the lemma below. To formulate it we need some more definitions.

Given an ordered  $k$ -graph  $G$ , let  $\mathcal{R}_G(n, d)$  be the family of *extensions* of  $G$ , that is, ordered  $d$ -regular  $k$ -graphs the first edges of which are equal to  $G$ . More precisely, setting  $G = (e_1, \dots, e_t)$ ,

$$\mathcal{R}_G(n, d) = \{H = (f_1, \dots, f_M) : f_i = e_i, i = 1, \dots, t, \text{ and } H \in \mathcal{R}^{(k)}(n, d)\}.$$

We say that a  $k$ -graph  $G$  with  $t \leq M$  edges is *admissible*, if  $\mathcal{R}_G(n, d) \neq \emptyset$  or, equivalently,  $\mathbb{P}(\mathbb{R}(t) = G) > 0$ . We define

$$p_{t+1}(e|G) := \mathbb{P}(\eta_{t+1} = e \mid \mathbb{R}(t) = G), \quad t = 0, \dots, M - 1. \quad (4)$$

Given  $\epsilon \in (0, 1)$ , we define events

$$\mathcal{A}_t = \left\{ p_{t+1}(e|\mathbb{R}(t)) \geq \frac{1 - \epsilon}{\binom{n}{k} - t} \quad \text{for every } e \in K_n \setminus \mathbb{R}(t) \right\}, \quad t = 0, \dots, M - 1. \quad (5)$$

Now we are ready to state the main ingredient of the proof of Theorem 1.

**Lemma 6.** *For every  $k \geq 2$  there is a positive constant  $C'$  such that if, for some real  $\epsilon = \epsilon(n)$  and integer  $d = d(n)$ ,*

$$C' \left( (d/n^{k-1} + (\log n)/d)^{1/3} + 1/n \right) \leq \epsilon < 1 \quad (6)$$

*and  $(1 - \epsilon)M$  is an integer, then the event  $\mathcal{A} := \mathcal{A}_0 \cap \dots \cap \mathcal{A}_{(1-\epsilon)M-1}$  occurs a.a.s.*

From Lemma 6, which is proved in Section 4, we deduce Theorem 1 using a coupling similar to the one which was used in [7].

*Proof of Theorem 1.* Clearly, we can pick  $\epsilon \leq \gamma/3$  such that  $(1 - \epsilon)M$  is integer and (1) implies (6) with  $C'$  being some constant multiple of  $C$ .

Let us first outline the proof. We will define a  $k$ -graph process  $\mathbb{R}'(t) := (\eta'_1, \dots, \eta'_t), t = 0, \dots, M$  such that for every admissible  $k$ -graph  $G$  with  $t \leq M - 1$  edges,

$$\mathbb{P}(\eta'_{t+1} = e \mid \mathbb{R}'(t) = G) = p_{t+1}(e \mid G). \quad (7)$$

In view of (7), the distribution of  $\mathbb{R}'(M)$  is the same as the one of  $\mathbb{R}(M)$  and thus we can define  $\mathbb{R}(n, d)$  as the  $k$ -graph  $\mathbb{R}'(M)$  with order of edges ignored. Then we will show that a.a.s.  $\mathbb{G}(n, m)$  can be sampled from the subhypergraph  $\mathbb{R}'((1 - \epsilon)M)$  of  $\mathbb{R}'(M)$ .

Now come the details. Set  $\mathbb{R}'(0)$  to be an empty vector and define  $\mathbb{R}'(t)$  inductively (for  $t = 1, 2, \dots$ ) as follows. Suppose that  $k$ -graphs  $R_t = \mathbb{R}'(t)$  and  $G_t = \mathbb{G}(t)$  have been exposed. Draw  $\varepsilon_{t+1}$  uniformly at random from  $K_n \setminus G_t$  and, independently, generate a Bernoulli random variable  $\xi_{t+1}$  with the probability of success  $1 - \epsilon$ . If event  $\mathcal{A}_t$  has occurred, that is,

$$p_{t+1}(e \mid R_t) \geq \frac{1 - \epsilon}{\binom{n}{k} - t} \quad \text{for every } e \in K_n \setminus R_t, \quad (8)$$

then draw a random edge  $\zeta_{t+1} \in K_n \setminus R_t$  according to the distribution

$$\mathbb{P}(\zeta_{t+1} = e \mid \mathbb{R}'(t) = R_t) := \frac{p_{t+1}(e \mid R_t) - (1 - \epsilon) / (\binom{n}{k} - t)}{\epsilon} \geq 0,$$

where the inequality holds by (8). Observe also that

$$\sum_{e \in K_n \setminus R_t} \mathbb{P}(\zeta_{t+1} = e \mid \mathbb{R}'(t) = R_t) = 1,$$

so  $\zeta_{t+1}$  has a well-defined distribution. Finally, fix an arbitrary bijection  $f_{R_t, G_t} : R_t \setminus G_t \rightarrow G_t \setminus R_t$  between the sets of edges and define

$$\eta'_{t+1} = \begin{cases} \varepsilon_{t+1}, & \text{if } \xi_{t+1} = 1, \varepsilon_{t+1} \in K_n \setminus R_t, \\ f_{R_t, G_t}(\varepsilon_{t+1}), & \text{if } \xi_{t+1} = 1, \varepsilon_{t+1} \in R_t, \\ \zeta_{t+1}, & \text{if } \xi_{t+1} = 0. \end{cases} \quad (9)$$

If the event  $\mathcal{A}_t$  fails, we nevertheless generate  $\xi_{t+1}$ , whereas  $\eta'_{t+1}$  is then sampled directly (without defining  $\zeta_{t+1}$ ) according to probabilities (4). Such a definition of  $\eta'_{t+1}$  ensures that

$$\mathcal{A}_t \cap \{\xi_{t+1} = 1\} \implies \varepsilon_{t+1} \in \mathbb{R}'(t+1). \quad (10)$$

Further, define a random subsequence of edges of  $\mathbb{G}((1 - \epsilon)M)$ ,

$$S := \{\varepsilon_i : \xi_i = 1, i \leq (1 - \epsilon)M\}.$$

Conditioning on the vector  $(\xi_i)$  determines  $|S|$ . If  $|S| \geq m$ , we define  $\mathbb{G}(n, m)$  as the first  $m$  edges of  $S$  (note that since the vectors  $(\xi_i)$  and  $(\varepsilon_i)$  are independent, these  $m$  edges are uniformly distributed), and if  $|S| < m$ , then we define  $\mathbb{G}(n, m)$  as a graph with edges  $\{\varepsilon_1, \dots, \varepsilon_m\}$ .

Let event  $\mathcal{A}$  be as in Lemma 6. The crucial thing is that by (10) we have

$$\mathcal{A} \implies S \subset \mathbb{R}'(M).$$

Therefore

$$\mathbb{P}(\mathbb{G}(n, m) \subset \mathbb{R}(n, d)) \geq \mathbb{P}(\{|S| \geq m\} \cap \mathcal{A}).$$

Since by Lemma 6 event  $\mathcal{A}$  holds a.a.s., to complete the proof it suffices to show that  $\mathbb{P}(|S| < m) \rightarrow 0$ .

To this end, note that  $|S|$  is a binomial random variable, namely,

$$|S| = \sum_{i=1}^{(1-\epsilon)M} \xi_i \sim \text{Bin}((1-\epsilon)M, 1-\epsilon),$$

with

$$\mathbb{E}|S| \geq (1-2\epsilon)M \quad \text{and} \quad \text{Var}|S| = (1-\epsilon)^2\epsilon M \leq \epsilon M. \quad (11)$$

Recall that  $\epsilon \leq \gamma/3$  and thus  $m = (1-\gamma)M \leq (1-3\epsilon)M$ . By (11), Chebyshev's inequality, and the inequality  $\epsilon \geq C' \log n/d$ , which follows from (6), we get

$$\mathbb{P}(|S| < m) \leq \mathbb{P}(|S| - \mathbb{E}|S| < -\epsilon M) \leq \frac{\epsilon M}{(\epsilon M)^2} = \frac{k}{\epsilon n d} \leq \frac{k}{C' n \log n} \rightarrow 0. \quad (12)$$

□

### 3 Preparations for the Proof of Lemma 6

Throughout this section we adopt the assumptions of Lemma 6, that is,  $(1-\epsilon)M$  is an integer and (6) holds with a sufficiently large  $C' = C'(k) \geq 1$ . In particular,

$$\epsilon \geq C'(\log n/d)^\alpha, \quad (13)$$

$$\epsilon \geq C'(d/n^{k-1})^\alpha \quad (14)$$

for every  $\alpha \geq 1/3$ , and

$$\epsilon \geq C'/n. \quad (15)$$

Also, let

$$\tau = 1 - \frac{t}{M}.$$



Given a  $k$ -graph  $G$  with maximum degree at most  $d$ , let us define the *residual degree* of a vertex  $v \in V(G)$  as

$$r_G(v) = d - \deg_G(v).$$

We begin our preparations toward the proof of Lemma 6 with a fact which allows one to control the residual degrees of the evolving  $k$ -graph  $\mathbb{R}(t) = (\eta_1, \dots, \eta_t)$ . For a vertex  $v \in [n]$  and  $t = 0, \dots, M$ , define random variables

$$R_t(v) = r_{\mathbb{R}(t)}(v) = |\{i \in (t, M] : v \in \eta_i\}|.$$

**Claim 7.** *For every  $k \geq 2$  there is a constant  $a = a(k) > 0$  such that a.a.s.*

$$\forall t \leq (1 - \epsilon)M, \quad \forall v \in [n], \quad |R_t(v) - \tau d| \leq \sqrt{a\tau d \log n} \leq \tau d/2 - 1. \quad (16)$$

*Proof.* A crucial observation is that the concentration of the degrees depends solely on the random ordering of the edges and not on the structure of the  $k$ -graph  $\mathbb{R}(M)$ . If we fix a  $d$ -regular  $k$ -graph  $H$  and condition  $\mathbb{R}(M)$  to be a random permutation of the edges of  $H$ , then  $R_t(v)$  is a hypergeometric random variable with expectation

$$\mathbb{E}R_t(v) = \frac{(M - t)d}{M} = \tau d,$$

and variance

$$\text{Var}R_t(v) = \frac{td(M - t)(M - d)}{M^2(M - 1)} \leq \frac{d(M - t)}{M} = \tau d.$$

Using Remark 2.11 in [9] together with inequalities (2.14) and (2.16) therein, we get

$$\mathbb{P}(|R_t(v) - \tau d| \geq x) \leq 2 \exp \left\{ -\frac{x^2}{2(\text{Var}R_t(v) + x/3)} \right\} \leq 2 \exp \left\{ -\frac{x^2}{2\tau d(1 + x/(3\tau d))} \right\}.$$

Let  $a = 3(k + 2)$  and  $x = \sqrt{a\tau d \log n}$ . Condition (13) with  $\alpha = 1$  and  $C' \geq 9a$  implies that

$$\tau d \geq \epsilon d \geq C' \log n. \quad (17)$$

Therefore

$$x/(\tau d) = \sqrt{a \log n / (\tau d)} \leq \sqrt{a \log n / (\epsilon d)} \leq \sqrt{a/C'} \leq 1/3. \quad (18)$$

Hence,

$$\mathbb{P}\left(|R_t(v) - \tau d| \geq \sqrt{a\tau d \log n}\right) \leq 2 \exp \left\{ -\frac{a}{3} \log n \right\} = 2n^{-k-2}.$$

Since we have fewer than  $nM \leq n^{k+1}$  choices of  $t$  and  $v$ , the first bound in (16) follows by taking the union bound.

The ultimate bound in (16) follows from (18), since

$$\sqrt{a\tau d \log n} = x \leq \tau d/3 \leq \tau d/2 - 1,$$

where the last inequality holds (for large enough  $n$ ) by (17).  $\square$

Recall that  $\mathcal{R}_G(n, d)$  is the family of extensions of  $G$  to a  $d$ -regular ordered  $k$ -graph. For a  $k$ -graph  $H \in \mathcal{R}_G(n, d)$  define the *common degree* (relative to subhypergraph  $G \subseteq H$ ) of an ordered pair  $(u, v)$  of vertices as

$$\text{cod}_{H|G}(u, v) = \left| \left\{ W \in \binom{[n]}{k-1} : W \cup u \in H, W \cup v \in H \setminus G \right\} \right|.$$

Note that  $\text{cod}_{H|G}(u, v)$  is not symmetric in  $u$  and  $v$ . Also, define the *degree of a pair of vertices*  $u, v$  as

$$\text{deg}_H(uv) = |\{e \in H : \{u, v\} \subset e\}|.$$

**Claim 8.** *Let  $G$  be an admissible  $k$ -graph with  $t + 1 \leq (1 - \epsilon)M$  edges such that*

$$r_G(v) \leq 2\tau d \quad \forall v \in [n]. \quad (19)$$

*Suppose that  $\mathbb{R}_G$  is a  $k$ -graph chosen uniformly at random from  $\mathcal{R}_G(n, d)$ . There are constants  $C_0, C_1$ , and  $C_2$ , depending on  $k$  only such that the following holds.*

*For each  $e \in K_n \setminus G$ ,*

$$\mathbb{P}(e \in \mathbb{R}_G) \leq \frac{C_0 \tau d}{n^{k-1}}. \quad (20)$$

*Moreover, if  $\ell \geq \ell_1 := C_1 \tau d/n$ , then for every  $u, v \in [n], u \neq v$ ,*

$$\mathbb{P}(\text{deg}_{\mathbb{R}_G \setminus G}(uv) > s) \leq 2^{-(\ell - \ell_1)}. \quad (21)$$

*Also, if  $\ell \geq \ell_2 := C_2 \tau d^2/n^{k-1}$ , then for every  $u, v \in [n], u \neq v$ ,*

$$\mathbb{P}(\text{cod}_{\mathbb{R}_G|G}(u, v) > \ell) \leq 2^{-(\ell - \ell_2)}. \quad (22)$$

*Proof.* To prove (20) define families of ordered  $k$ -graphs

$$\mathcal{R}_{e \in} = \{H \in \mathcal{R}_G(n, d) : e \in H\} \quad \text{and} \quad \mathcal{R}_{e \notin} = \{H \in \mathcal{R}_G(n, d) : e \notin H\}.$$

and observe that

$$\mathbb{P}(e \in \mathbb{R}_G) \leq \frac{|\mathcal{R}_{e \in}|}{|\mathcal{R}_{e \notin}|}.$$

In order to estimate this ratio, define an auxiliary bipartite graph  $B$  between  $\mathcal{R}_{e \in}$  and  $\mathcal{R}_{e \notin}$  in which  $H \in \mathcal{R}_{e \in}$  is connected to  $H' \in \mathcal{R}_{e \notin}$  whenever  $H'$  can be obtained from  $H$  by the following operation (known as *switching* in the literature dating back to McKay [13]). Let  $e = e_1 = \{v_{1,1} \dots v_{1,k}\}$  and pick  $k - 1$  more edges

$$e_i = \{v_{i,1} \dots v_{i,k}\} \in H \setminus G, \quad i = 2, \dots, k$$

(with vertices in the increasing order within each edge) so that all  $k$  edges are disjoint. Replace, for each  $j = 1, \dots, k$ , the edge  $e_j$  by

$$f_j := \{v_{1,j} \dots v_{k,j}\}$$

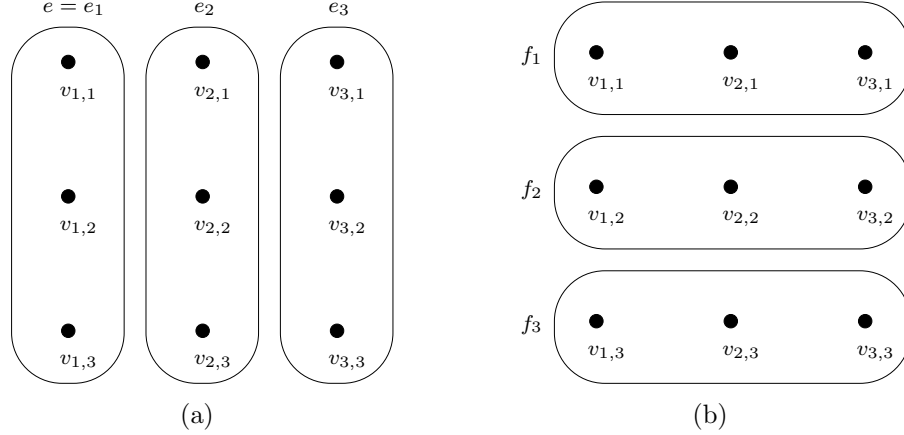


Figure 1: Switching (for  $k = 3$ ): before (a) and after (b).

to obtain  $H'$  (see Figure 1).

Let  $f(H)$  be the number of  $k$ -graphs  $H' \in \mathcal{R}_{e \notin}$  which can be obtained from  $H$ , and  $b(H')$  be the number of  $k$ -graphs  $H \in \mathcal{R}_{e \in}$  from which  $H'$  can be obtained. Thus,

$$|\mathcal{R}_{e \in}| \cdot \min_{H \in \mathcal{R}_{e \in}} f(H) \leq |E(B)| \leq |\mathcal{R}_{e \notin}| \cdot \max_{H'} b(H'). \quad (23)$$

Note that  $H \setminus G$  and  $H' \setminus G$  each have  $\tau M - 1$  edges and, by (19), maximum degrees at most  $2\tau d$ . To estimate  $f(H)$ , note that because each edge intersects at most  $k \cdot 2\tau d$  other edges of  $H \setminus G$ , the number of ways to choose an unordered  $(k-1)$ -tuple  $\{e_2, \dots, e_k\}$  is at least

$$\frac{1}{(k-1)!} \prod_{i=1}^{k-1} (\tau M - 1 - ik \cdot 2\tau d) \geq (\tau M - k^2 \cdot 2\tau d)^{k-1} / (k-1)!. \quad (24)$$

The number of such  $(k-1)$ -tuples that may lead to a double edge after the switching (by repeating some edge of  $H$  which intersects  $e_1$ ), is at most  $kd \cdot (2\tau d)^{k-1}$ . Thus,

$$\begin{aligned} f(H) &\geq \frac{(\tau M - 2k^2\tau d)^{k-1}}{(k-1)!} - k(2\tau)^{k-1}d^k \\ &= \frac{(\tau M)^{k-1}}{(k-1)!} \left( \left(1 - \frac{2k^2d}{M}\right)^{k-1} - \frac{k!(2\tau)^{k-1}d^k}{(\tau M)^{k-1}} \right) \\ &= \frac{(\tau M)^{k-1}}{(k-1)!} \left( \left(1 - \frac{2k^3}{n}\right)^{k-1} - \frac{k!(2k)^{k-1}d}{n^{k-1}} \right) \\ &\geq \frac{(\tau M)^{k-1}}{(k-1)!} \left( 1 - \frac{2k^4}{n} - \frac{(2k)^{2k}d}{n^{k-1}} \right). \end{aligned}$$

By (14) with  $\alpha = 1$ , (15), and sufficiently large  $C'$ , we have

$$\frac{2k^4}{n} + \frac{(2k)^{2k}d}{n^{k-1}} \leq \frac{\epsilon(2k^4 + (2k)^{2k})}{C'} \leq 1/2.$$

Hence,

$$f(H) \geq \frac{(\tau M)^{k-1}}{2(k-1)!}. \quad (25)$$

In order to bound  $b(H')$  from above note that there are at most  $(2\tau d)^k$  ways to choose a sequence  $f_1, \dots, f_k \in H' \setminus G$  such that  $v_{1,i} \in f_i$  and we can reconstruct the  $k-1$ -tuple  $e_2, \dots, e_k$  in at most  $((k-1)!)^{k-1}$  ways (by fixing an ordering of vertices of  $f_1$  and permuting vertices in other  $f_i$ 's). Therefore  $b(H') \leq ((k-1)!)^{k-1} \cdot (2\tau d)^k$ . This, with (23) and (25) implies that

$$\mathbb{P}(e \in \mathbb{R}_G) \leq \frac{|\mathcal{R}_{e \in}|}{|\mathcal{R}_{e \notin}|} \leq \frac{\max_{H' \in \mathcal{R}_{e \notin}} b(H')}{\min_{H \in \mathcal{R}_{e \in}} f(H)} \leq \frac{2((k-1)!)^k (2\tau d)^k}{(\tau M)^{k-1}} = \frac{C_0 \tau d}{n^{k-1}},$$

for some constant  $C_0 = C_0(k)$ . This concludes the proof of (20).

To prove (21), fix  $u, v \in [n]$  and define the families

$$\mathcal{R}_1(\ell) = \{H \in \mathcal{R}_G(n, d) : \deg_{H \setminus G}(uv) = \ell\}, \quad \ell = 0, 1, \dots$$

In order to compare sizes of  $\mathcal{R}_1(\ell)$  and  $\mathcal{R}_1(\ell-1)$  we define the following switching which maps a  $k$ -graph  $H \in \mathcal{R}_1(\ell)$  to a  $k$ -graph  $H' \in \mathcal{R}_1(\ell-1)$ . Select  $e_1 \in H \setminus G$  contributing to  $\deg_{H \setminus G}(uv)$  and pick  $k-1$  edges  $e_2, \dots, e_k \in H \setminus G$  so that  $e_1, \dots, e_k$  are disjoint. Writing  $e_i = v_{i,1} \dots v_{i,k}$ ,  $i = 1, \dots, k$  with  $u = v_{1,1}$  and  $v = v_{1,2}$ , replace  $e_1, \dots, e_k$  by  $f_j = v_{1,j} \dots v_{k,j}$ ,  $j = 1, \dots, k$  (as in Figure 1).

Noting that this time  $e_1$  can be chosen in  $\ell$  ways, we get a lower bound on  $f(H)$  very similar to that in (25):

$$f(H) \geq \ell \left( (\tau M - 2k^2 \tau d)^{k-1} / (k-1)! - k(2\tau)^{k-1} d^k \right) \geq \frac{\ell(\tau M)^{k-1}}{2(k-1)!}.$$

For the upper bound for  $b(H')$  we choose two disjoint edges in  $H' \setminus G$  containing  $u$  and  $v$ , respectively, and then  $k-2$  more edges in  $H' \setminus G$  not containing  $u$  and  $v$  so that all edges are disjoint. Crudely bounding number of permutations of vertices inside each of  $f_1, \dots, f_k$  by  $(k!)^k$ , we get  $b(H') \leq (k!)^k (2\tau d)^2 (\tau M)^{k-2}$ . We obtain

$$\frac{|\mathcal{R}_1(\ell)|}{|\mathcal{R}_1(\ell-1)|} \leq \frac{\max_{H' \in \mathcal{R}_1(\ell-1)} b(H')}{\min_{H \in \mathcal{R}_1(\ell)} f(H)} \leq \frac{2(k!)^{k+1} (2\tau d)^2 (\tau M)^{k-2}}{\ell(\tau M)^{k-1}} \leq \frac{8(k!)^{k+1} \tau d}{\ell n} \leq \frac{1}{2},$$

by assumption  $\ell \geq \ell_1 = C_1 \tau d/n$  and appropriate choice of constant  $C_1$ . Further,

$$\begin{aligned} \mathbb{P}(\deg_{\mathbb{R}_G \setminus G}(u, v) > \ell) &\leq \sum_{i > \ell} \frac{|\mathcal{R}_1(i)|}{|\mathcal{R}_G(n, d)|} \leq \sum_{i > \ell} \frac{|\mathcal{R}_1(i)|}{|\mathcal{R}_1(\ell_1)|} \\ &= \sum_{i > \ell} \prod_{j=\ell_1+1}^i \frac{|\mathcal{R}_1(j)|}{|\mathcal{R}_1(j-1)|} \leq \sum_{i > \ell} 2^{-(i-\ell_1)} = 2^{-(\ell-\ell_1)}, \end{aligned} \quad (26)$$

which completes the proof of (21).

It remains to show (22). Fix  $u, v \in [n]$  and define the families

$$\mathcal{R}_2(\ell) = \{H \in \mathcal{R}_G(n, d) : \text{cod}_{H|G}(u, v) = \ell\}, \quad \ell = 0, 1, \dots$$

We compare sizes of  $\mathcal{R}_2(\ell)$  and  $\mathcal{R}_2(\ell-1)$  using the following switching. Select two edges  $e_0 \in H$  and  $e_1 \in H \setminus G$  contributing to  $\text{cod}_{H|G}(u, v)$ , that is, such that  $e_0 \setminus u = e_1 \setminus v$ ; pick  $k-1$  other edges  $e_2, \dots, e_k \in H \setminus G$  so that  $e_1, \dots, e_k$  are disjoint. Writing  $e_i = v_{i,1} \dots v_{i,k}$ ,  $i = 1, \dots, k$  with  $v = v_{1,1}$ , replace  $e_1, \dots, e_k$  by  $f_j = v_{1,j} \dots v_{k,j}$ ,  $j = 1, \dots, k$  (see Figure 2).

We estimate  $f(H)$  by first fixing a pair  $e_0, e_1$  in one of  $\ell$  ways. The number of choices of  $e_2, \dots, e_k$  is bounded as in (24). However, we subtract not just at most  $kd \cdot (2\tau d)^{k-1}$   $(k-1)$ -tuples which may create double edges, but also  $(k-1)$ -tuples for which  $f_1 \setminus \{v\} \cup \{u\} \in H$  which prevents  $\text{cod}(u, v)$  from being decreased. There are at most  $d \cdot (2\tau d)^{k-1}$  of such  $(k-1)$ -tuples, hence

$$\begin{aligned} f(H) &\geq \ell \left( \frac{(\tau M - k^2 \cdot 2\tau d)^{k-1}}{(k-1)!} - (k+1)d \cdot (2\tau d)^{k-1} \right) \\ &= \frac{\ell(\tau M)^{k-1}}{(k-1)!} \left( \left(1 - \frac{2k^2 d}{M}\right)^{k-1} - (k+1)(k-1)!d \left(\frac{2d}{M}\right)^{k-1} \right) \\ &= \frac{\ell(\tau M)^{k-1}}{(k-1)!} \left( \left(1 - \frac{2k^3}{n}\right)^{k-1} - (k+1)(k-1)!d \left(\frac{2k}{n}\right)^{k-1} \right) \\ &\geq \frac{\ell(\tau M)^{k-1}}{(k-1)!} \left( 1 - \frac{2k^4}{n} - (k+1)!(2k)^k \frac{d}{n^{k-1}} \right) \\ &\geq \frac{\ell(\tau M)^{k-1}}{2(k-1)!}, \end{aligned}$$

where the last inequality follows from (14) with  $\alpha = 1$  and (15) with sufficiently large  $C'$ .

Conversely,  $H$  can be reconstructed from  $H'$  by choosing an edge  $e_0 \in H'$  containing  $u$  but not containing  $v$  and then  $k$  disjoint edges  $f_j \in H' \setminus G$ , each containing

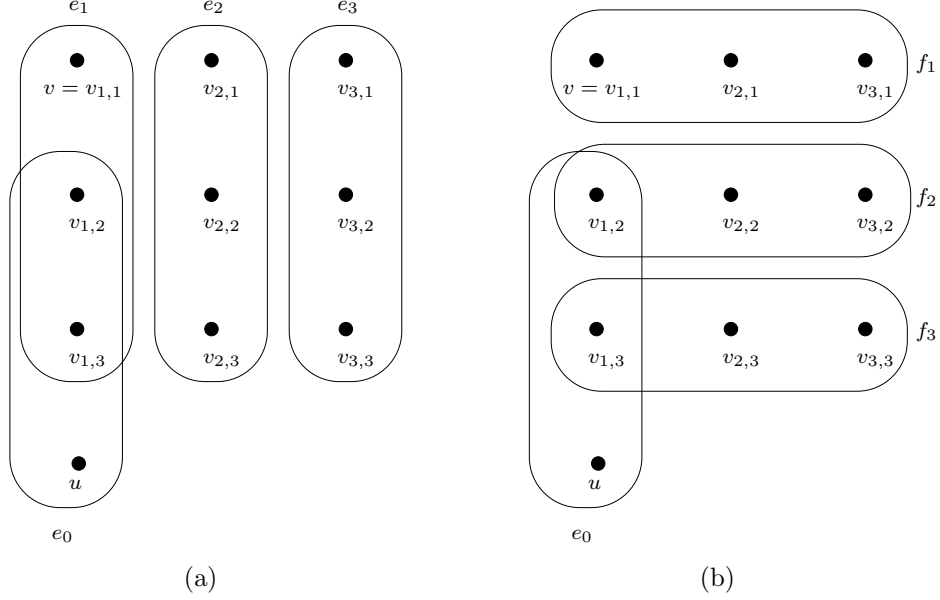


Figure 2: Switching (for  $k = 3$ ): before (a) and after (b).

exactly one vertex from  $(e_0 \setminus u) \cup v$  and permuting the vertices inside  $f_2 \setminus v_{1,2}, \dots, f_k \setminus v_{1,k}$  in at most  $((k-1)!)^{k-1}$  ways. Therefore  $b(H') \leq ((k-1)!)^{k-1} d(2\tau d)^k$ . Clearly,

$$\frac{|\mathcal{R}_2(\ell)|}{|\mathcal{R}_2(\ell-1)|} \leq \frac{\max_{H' \in \mathcal{R}_2(\ell-1)} b(H')}{\min_{H \in \mathcal{R}_2(\ell)} f(H)} \leq \frac{d(2\tau d)^k \cdot 2((k-1)!)^k}{\ell(\tau M)^{k-1}} \leq \frac{2^{k+1}((k-1)!)^k k^{k-1} \tau d^2}{n^{k-1} \ell} \leq \frac{1}{2},$$

by the assumption  $\ell \geq \ell_2 = C_2 \tau d^2 / n^{k-1}$  and appropriate choice of constant  $C_2$ . Now (22) follows from similar computations to (21).

This finishes the proof of Claim 8.  $\square$

## 4 Proof of Lemma 6

In this section we prove the crucial Lemma 6. In view of Claim 7 it suffices to show that

$$\mathbb{P}(\eta_{t+1} = e \mid \mathbb{R}(t) = G) \geq \frac{1-\epsilon}{\binom{n}{k} - t}, \quad \forall e \in K_n \setminus G, \quad (27)$$

for every  $t \leq (1-\epsilon)M - 1$  and every admissible  $G$  such that

$$d(\tau - \delta) \leq r_G(v) \leq d(\tau + \delta), \quad v \in [n], \quad (28)$$

where

$$\tau = 1 - t/M \quad \text{and} \quad \delta = \sqrt{a\tau(\log n)/d}.$$

In some cases the following simpler bounds (implied by the second inequality in (16)) on  $r_G(v)$  will suffice:

$$\tau d/2 + 1 \leq r_G(v) \leq 2\tau d, \quad v \in [n]. \quad (29)$$

Since the average of  $\mathbb{P}(\eta_{t+1} = e \mid \mathbb{R}(t) = G)$  over  $e \in K_n \setminus G$  is exactly  $1 / \binom{n}{k} - t$ , there is  $f \in K_n \setminus G$  such that

$$\mathbb{P}(\eta_{t+1} = f \mid \mathbb{R}(t) = G) \geq \frac{1}{\binom{n}{k} - t}. \quad (30)$$

Fix any such  $f$  and let  $e \in K_n \setminus G$  be arbitrary. Setting  $\mathcal{R}_f := \mathcal{R}_{G \cup f}(n, d)$  and  $\mathcal{R}_e := \mathcal{R}_{G \cup e}(n, d)$ , we have

$$\frac{\mathbb{P}(\eta_{t+1} = e \mid \mathbb{R}(t) = G)}{\mathbb{P}(\eta_{t+1} = f \mid \mathbb{R}(t) = G)} = \frac{|\mathcal{R}_{G \cup e}(n, d)|}{|\mathcal{R}_{G \cup f}(n, d)|} = \frac{|\mathcal{R}_e|}{|\mathcal{R}_f|}. \quad (31)$$

To bound this ratio, we need to appeal to the configuration model for hypergraphs. Let  $\mathbb{M}_G(n, d)$  be a random *multi- $k$ -graph extension* of  $G$  to an ordered  $d$ -regular multi- $k$ -graph. Namely,  $\mathbb{M}_G(n, d)$  is a sequence of  $M$  edges (each of which is a  $k$ -element multiset of vertices), the first  $t$  of which comprise  $G$ , while the remaining ones are generated by taking a random uniform permutation  $\Pi$  of the multiset

$$\{1, \dots, 1, \dots, n, \dots, n\}$$

with multiplicities  $r_G(v)$ ,  $v \in [n]$ , and splitting it into consecutive  $k$ -tuples.

The number of such permutations is

$$N_G := \frac{(k(M-t))!}{\prod_{v \in [n]} r_G(v)!}.$$

Since each simple extension of  $G$  is given by the same number  $(k!)^{M-t}$  of permutations,  $\mathbb{M}_G(n, d)$  is uniform over  $\mathcal{R}_G(n, d)$ . That is,  $\mathbb{M}_G(n, d)$ , conditioned on simplicity, has the same distribution as  $\mathbb{R}_G(n, d)$ .

Set

$$\mathbb{M}_e = \mathbb{M}_{G \cup e}(n, d) \quad \text{and} \quad \mathbb{M}_f = \mathbb{M}_{G \cup f}(n, d),$$

for convenience. Noting that  $G \cup f$  has  $t+1$  edges, we have

$$\mathbb{P}(\mathbb{M}_f \in \mathcal{R}_f) = \frac{|\mathcal{R}_f| (k!)^{M-t-1}}{N_{G \cup f}} = \frac{|\mathcal{R}_f| (k!)^{M-t-1} \prod_{v \in [n]} r_{G \cup f}(v)!}{(k(M-t-1))!},$$

and similarly for  $\mathbb{M}_e$  and  $\mathcal{R}_e$ . This yields, after a few cancelations, that

$$\frac{|\mathcal{R}_e|}{|\mathcal{R}_f|} = \frac{\prod_{v \in e \setminus f} r_G(v)}{\prod_{v \in f \setminus e} r_G(v)} \cdot \frac{\mathbb{P}(\mathbb{M}_e \in \mathcal{R}_e)}{\mathbb{P}(\mathbb{M}_f \in \mathcal{R}_f)}. \quad (32)$$

The ratio of the products in (32) is, by (28), at least

$$\left(\frac{\tau - \delta}{\tau + \delta}\right)^k \geq \left(1 - \frac{2\delta}{\tau}\right)^k \geq 1 - 2k\sqrt{\frac{a \log n}{\tau d}} \geq 1 - 2k\sqrt{\frac{a \log n}{\epsilon d}} \geq 1 - \epsilon/2,$$

where the last inequality holds by (13) with  $\alpha = 1/3$  and  $C' \geq \sqrt[3]{16ak^2}$ . On the other hand, the ratio of probabilities in (32) will be shown in Claim 9 below to be at least  $1 - \epsilon/2$ . Consequently, the entire ratio in (32), and thus in (31), will be at least  $1 - \epsilon$ , which, in view of (30), will imply (27) and yield the lemma.

Hence, to complete the proof of Lemma 6 it remains to show that the probabilities of simplicity  $\mathbb{P}(\mathbb{M}_e \in \mathcal{R}_e)$  are asymptotically the same for all  $e \in K_n \setminus G$ . Recall that for every edge  $e \in K_n \setminus G$  we write

$$\mathbb{M}_e = \mathbb{M}_{G \cup e}(n, d) \quad \text{and} \quad \mathcal{R}_e = \mathcal{R}_{G \cup e}(n, d). \quad (33)$$

**Claim 9.** *If  $G$ ,  $e$ , and  $f$  are as above, then, for every  $e \in K_n \setminus G$ ,*

$$\frac{\mathbb{P}(\mathbb{M}_e \in \mathcal{R}_e)}{\mathbb{P}(\mathbb{M}_f \in \mathcal{R}_f)} \geq 1 - \epsilon/2.$$

*Proof.* We start by constructing a coupling of  $\mathbb{M}_e$  and  $\mathbb{M}_f$  in which they differ in at most  $k + 1$  edges (counting in the replacement of  $f$  by  $e$  at the  $(t + 1)$ -th position).

Let  $f = u_1 \dots u_k$  and  $e = v_1 \dots v_k$ . Further, let  $r = k - |f \cap e|$  and suppose without loss of generality that  $\{u_1 \dots u_r\} \cap \{v_1 \dots v_r\} = \emptyset$ . Let  $\Pi_f$  be a random permutation underlying the multi- $k$ -graph  $\mathbb{M}_f$ . Note that  $\Pi_f$  differs from any permutation  $\Pi_e$  underlying  $\mathbb{M}_e$  by having the multiplicities of  $v_1, \dots, v_r$  greater by one, and the multiplicities of  $u_1, \dots, u_r$  smaller by one than the corresponding multiplicities in  $\Pi_e$ .

Let  $\Pi^*$  be obtained from  $\Pi_f$  by replacing, for each  $i = 1, \dots, r$ , a copy of  $v_i$  selected uniformly at random by  $u_i$ . Define  $\mathbb{M}^*$  by chopping  $\Pi^*$  into consecutive  $k$ -tuples and appending them to  $G \cup e$  (see Figure 3).

It is easy to see that  $\Pi^*$  is uniform over all permutations of the multiset

$$\{1, \dots, 1, \dots, n, \dots, n\}$$

with multiplicities  $r_{G \cup e}(v)$ ,  $v \in [n]$ . This means that  $\mathbb{M}^*$  has the same distribution as  $\mathbb{M}_e$  and thus we will further identify  $\mathbb{M}^*$  and  $\mathbb{M}_e$ .

Observe that if we condition  $\mathbb{M}_f$  on being a simple  $k$ -graph  $H$ , then  $\mathbb{M}_e$  can be equivalently obtained by the following switching: (i) replace edge  $f$  by  $e$ ; (ii) for each  $i = 1, \dots, r$ , choose, uniformly at random, an edge  $e_i \in H \setminus (G \cup f)$  incident to  $v_i$  and replace it by  $(e_i \setminus v_i) \cup u_i$  (see Figure 4). Of course, some of  $e_i$ 's may coincide. For example, if  $e_{i_1} = \dots = e_{i_l}$ , then the effect of the switching is that  $e_{i_1}$  is replaced by  $(e_{i_1} \setminus \{v_{i_1}, \dots, v_{i_l}\}) \cup \{u_{i_1}, \dots, u_{i_l}\}$ .





The crucial idea is that such a switching is unlikely to create loops or multiple edges. However, for certain  $H$  this might not be true. For example, if  $e \in H \setminus (G \cup f)$ , then the random choice of  $e_i$ 's in step (ii) is unlikely to destroy  $e$ , but in step (i) edge  $f$  has been replaced by an additional copy of  $e$ , thus creating a double edge. Moreover, if almost every  $(k-1)$ -tuple of vertices extending  $v_i$  to an edge in  $H \setminus (G \cup f)$  also extends  $u_i$  to an edge in  $H$ , then most likely the replacement of  $v_i$  by  $u_i$  will create a double edge too. To avoid such and other bad instances, we say that  $H \in \mathcal{R}_f$  is *nice* if the following three properties hold

$$e \notin H \tag{34}$$

$$\max_{i=1, \dots, r} \deg_{H \setminus (G \cup f)}(u_i v_i) \leq \ell_1 + k \log_2 n, \tag{35}$$

$$\max_{i=1, \dots, r} \text{cod}_{H \setminus (G \cup f)}(u_i, v_i) \leq \ell_2 + k \log_2 n, \tag{36}$$

where  $\ell_1 = C_1 \tau d / n$  and  $\ell_2 = C_2 \tau d^2 / n^{k-1}$  are as in Claim 8. Note that  $\mathbb{M}_f$ , conditioned on  $\mathbb{M}_f \in \mathcal{R}_f$ , is distributed uniformly over  $\mathcal{R}_{G \cup f}(n, d)$ . Since we chose  $f$  such that by (30) is satisfied, we have that  $k$ -graph  $G \cup f$  is admissible. Therefore by Claim 8 we have

$$\begin{aligned} \mathbb{P}(\mathbb{M}_f \text{ is not nice} \mid \mathbb{M}_f \in \mathcal{R}_f) &\leq \frac{C_0 \tau d}{n^{k-1}} + 2 \cdot r 2^{-k \log_2 n} \\ &\leq \frac{C_0 d + 2k}{n^{k-1}} \leq \frac{\epsilon}{4}, \end{aligned} \tag{37}$$

where the last inequality follows by (14) with  $\alpha = 1$  and sufficiently large constant  $C'$ . By standard probability, we have

$$\begin{aligned} \frac{\mathbb{P}(\mathbb{M}_e \in \mathcal{R}_e)}{\mathbb{P}(\mathbb{M}_f \in \mathcal{R}_f)} &\geq \mathbb{P}(\mathbb{M}_e \in \mathcal{R}_e \mid \mathbb{M}_f \in \mathcal{R}_f) \\ &\geq \mathbb{P}(\mathbb{M}_e \in \mathcal{R}_e \mid \mathbb{M}_f \text{ is nice}) \mathbb{P}(\mathbb{M}_f \text{ is nice} \mid \mathbb{M}_f \in \mathcal{R}_f). \end{aligned} \tag{38}$$

It suffices to show that

$$\mathbb{P}(\mathbb{M}_e \in \mathcal{R}_e \mid \mathbb{M}_f \text{ is nice}) \geq 1 - \epsilon/4, \tag{39}$$

since in view of (37) and (39), inequality (38) completes the proof of the claim.

Now we prove (39). Fix a nice  $k$ -graph  $H \in \mathcal{R}_f$  and condition on the event  $\mathbb{M}_f = H$ . The event that  $\mathbb{M}_e$  is not simple is contained in the union of the following four events:

$$\mathcal{E}_1 = \{ \text{two of the randomly chosen edges } e_1, \dots, e_r \text{ coincide} \},$$

$$\mathcal{E}_2 = \{ (e_i \setminus v_i) \cup u_i \text{ is a loop for some } i = 1, \dots, r \},$$

$$\mathcal{E}_3 = \{ (e_i \setminus v_i) \cup u_i \in H \text{ for some } i = 1, \dots, r \},$$

$$\mathcal{E}_4 = \{ (e_i \setminus v_i) \cup u_i = (e_j \setminus v_j) \cup u_j \text{ for some distinct } i \text{ and } j \}.$$

Event  $\mathcal{E}_1$  covers all cases when a double edge is created by replacing several vertices in the same edge. Creation of multiple edges in other ways is addressed by events  $\mathcal{E}_3$  and  $\mathcal{E}_4$ .

In what follows we will several times use the fact that

$$\deg_{H \setminus (G \cup f)}(v) \geq \tau d / 2 \geq \epsilon d / 2, \quad \forall v \in [n], \quad (40)$$

which is immediate from (29) and  $\tau \geq \epsilon$ . To bound the probability of  $\mathcal{E}_1$ , observe that, given  $1 \leq i < j \leq r$ , the number of choices of a coinciding pair  $e_i = e_j$  is  $\deg_{H \setminus (G \cup f)}(v_i v_j) \leq \deg_{H \setminus (G \cup f)}(v_i)$  and the probability that both  $v_i$  and  $v_j$  actually select a fixed common edge is  $(\deg_{H \setminus (G \cup f)}(v_i) \deg_{H \setminus (G \cup f)}(v_j))^{-1}$ . Therefore using (40) we obtain

$$\begin{aligned} \mathbb{P}(\mathcal{E}_1 | \mathbb{M}_f = H) &\leq \sum_{1 \leq i < j \leq r} \frac{\deg_{H \setminus (G \cup f)}(v_i v_j)}{\deg_{H \setminus (G \cup f)}(v_i) \deg_{H \setminus (G \cup f)}(v_j)} \leq \sum_{1 \leq i < j \leq r} \frac{1}{\deg_{H \setminus (G \cup f)}(v_j)} \\ &\leq \frac{2 \binom{k}{2}}{\epsilon d} \leq \frac{\epsilon}{16}, \quad (41) \end{aligned}$$

where the last inequality follows from (13) with  $\alpha = 1/2$  and sufficiently large  $C'$ .

To bound the probability of  $\mathcal{E}_2$ , note that a loop in  $\mathbb{M}_e$  can only be created when for some  $i = 1, \dots, r$ , the randomly chosen edge  $e_i$  contains both  $v_i$  and  $u_i$ . There are at most  $\deg_{H \setminus (G \cup f)}(u_i v_i)$  such edges. Therefore, by (35) and (40) we get

$$\begin{aligned} \mathbb{P}(\mathcal{E}_2 | \mathbb{M}_f = H) &\leq \sum_{i=1}^r \frac{\deg_{H \setminus (G \cup f)}(u_i v_i)}{\deg_{H \setminus (G \cup f)}(v_i)} \leq \frac{2k(\ell_1 + k \log_2 n)}{\tau d} \\ &\leq \frac{2k\ell_1}{\tau d} + \frac{2k^2 \log_2 n}{\epsilon d} = \frac{2kC_1}{n} + \frac{2k^2 \log_2 n}{\epsilon d} \leq \frac{\epsilon}{16}, \quad (42) \end{aligned}$$

where the last inequality is implied by (13) with  $\alpha = 1/2$ , (15) and sufficiently large  $C'$ .

Similarly we bound the probability of  $\mathcal{E}_3$ , the event that for some  $i$  we will choose  $e_i \in H \setminus (G \cup f)$  with  $(e_i \setminus v_i) \cup u_i \in H$ . There are  $\text{cod}_{H|G \cup f}(u_i, v_i)$  such edges. Thus, by (36) and (40) we obtain

$$\begin{aligned} \mathbb{P}(\mathcal{E}_3 | \mathbb{M}_f = H) &\leq \sum_{i=1}^r \frac{\text{cod}_{H|G \cup f}(u_i, v_i)}{\deg_{H \setminus (G \cup f)}(v_i)} \leq \frac{2k(\ell_2 + k \log_2 n)}{\tau d} \\ &\leq \frac{2k\ell_2}{\tau d} + \frac{2k^2 \log_2 n}{\tau d} \leq \frac{2kC_2 d}{n^{k-1}} + \frac{2k \log_2 n}{\epsilon d} \leq \frac{\epsilon}{16}, \quad (43) \end{aligned}$$

where the last inequality follows from (13) with  $\alpha = 1/2$ , (14) with  $\alpha = 1$  and sufficiently large  $C'$ .

Finally, note that, given  $1 \leq i < j \leq r$ , if a pair  $e_i, e_j \in H \setminus (G \cup f)$  satisfies the condition in  $\mathcal{E}_4$ , then the edge  $e_j$  is uniquely determined by  $e_i$ . Therefore the number of such pairs is at most  $\deg_{H \setminus (G \cup f)}(v_i)$  and we get exactly the same bound as in (41):

$$\mathbb{P}(\mathcal{E}_4 \mid \mathbb{M}_f = H) \leq \sum_{1 \leq i < j \leq r} \frac{1}{\deg_{H \setminus (G \cup f)}(v_j)} \leq \frac{\epsilon}{16}. \quad (44)$$

Combining (41)-(44) and averaging over nice  $H$ , we obtain (39), as required.  $\square$

## 5 Concluding Remarks

Theorem 1 remains valid if we replace random hypergraph  $\mathbb{G}^{(k)}(n, m)$  by  $\mathbb{G}^{(k)}(n, p)$  with  $p = (1 - 2\gamma)d/\binom{n-1}{k-1}$ , say. To see this one can modify the proof of Theorem 1 as follows. Let  $B_n \sim \text{Bin}(\binom{n}{k}, p)$  be a random variable independent of the process  $(\mathbb{G}(t))_t$ . If  $B_n \leq m \leq |S|$ , sample  $\mathbb{G}^{(k)}(n, p)$  by taking the first  $B_n$  edges of  $S$  (which are uniformly distributed over all  $k$ -graphs with  $B_n$  edges). Otherwise sample  $\mathbb{G}^{(k)}(n, p)$  among  $k$ -graphs with  $B_n$  edges independently. In view of the assumption (3), Chernoff's inequality (see [9, (2.5)]) and (12) imply

$$\mathbb{P}(\mathbb{G}^{(k)}(n, p) \not\subset \mathbb{R}^{(k)}(n, d)) \leq \mathbb{P}(B_n > m) + \mathbb{P}(|S| < m) \rightarrow 0, \quad \text{as } n \rightarrow \infty.$$

The lower bound on  $d$  in Theorem 1 is necessary because the second moment method applied to  $\mathbb{G}^{(k)}(n, p)$  (cf. Theorem 3.1(ii) in [2]) and asymptotic equivalence of  $\mathbb{G}^{(k)}(n, p)$  and  $\mathbb{G}^{(k)}(n, m)$  yields that for  $d = o(\log n)$  and  $m \sim cM$  there is a sequence  $\Delta = \Delta(n) \gg d$  such that the maximum degree  $\mathbb{G}^{(k)}(n, m)$  is at least  $\Delta$  a.a.s.

In view of the above, our approach cannot be extended to  $d = O(\log n)$  in part (i) of Theorem 5. Nevertheless, we believe (as it was already stated in [7]) that for loose Hamilton cycles it suffices to assume that  $d = \Omega(1)$ .

**Conjecture 1.** *For every  $k \geq 3$  there is a constant  $d_k$  such that if  $d \geq d_k$ , then a.a.s.  $\mathbb{R}^{(k)}(n, d)$  contains a loose Hamilton cycle.*

We also believe that the lower bounds on  $d$  in parts (ii) and (iii) of Theorem 5 are of optimal order.

**Conjecture 2.** *For all integers  $k > \ell \geq 2$  if  $d \ll n^{\ell-1}$ , then a.a.s.  $\mathbb{R}^{(k)}(n, d)$  is not  $\ell$ -Hamiltonian.*

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