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RANDOM GRAPHS

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1 INTRODUCTION

Let $G = (V, E)$ be a graph in which the edges are coloured. A set $S \subseteq E$ is said to be *multicoloured* if each edge of S is a different colour. A spanning tree of G is said to be multicoloured if its edge set is. In this paper we study

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the existence of a multicoloured spanning tree (MST) in a randomly coloured random graph.

In fact, our main result will concern a randomly coloured graph process. Here e_1, e_2, \dots, e_N is a random permutation of the edges of the complete graph K_n and so $N = \binom{n}{2}$. Each edge e independently chooses a random colour $c(e)$ from a given set of colours W , $|W| \geq n - 1$.

The graph process consists of the sequence of random graphs $G_m, m = 1, 2, \dots, N$, where $G_m = ([n], E_m)$ and $E_m = \{e_1, e_2, \dots, e_m\}$. We identify the following events:

$$\mathcal{C}_m = \{G_m \text{ is connected}\}.$$

$$\mathcal{N}_m = \{|c(E_m)| \geq n - 1\}, \text{ where } c(E_m) \text{ is the set of colours used by } E_m.$$

$$\mathcal{MT}_m = \{G_m \text{ has a multicoloured spanning tree}\}.$$

Let \mathcal{E}_m stand for one of the above three sequences of events and let

$$m_{\mathcal{E}} = \min\{m : \mathcal{E}_m \text{ occurs}\},$$

provided such an m exists. Clearly, if $m_{\mathcal{MT}}$ is defined,

$$m_{\mathcal{MT}} \geq \max\{m_{\mathcal{C}}, m_{\mathcal{N}}\},$$

and the main result of the paper is

Theorem 1 *In almost every (a.e.) randomly coloured graph process*

$$m_{\mathcal{MT}} = \max\{m_{\mathcal{C}}, m_{\mathcal{N}}\}.$$

□

To establish the existence of an MST we use a result of Edmonds [2] on the matroid intersection problem. In this scenario M_1, M_2 are matroids over a common ground set E with rank functions r_1, r_2 respectively. Edmonds' general theorem on this problem is

$$\max(|I| : I \text{ is independent in both matroids}) = \min_{E_1 \cap E_2 = \emptyset} (r_1(E_1) + r_2(E_2)) \quad (1)$$

For us M_1 is the cycle matroid of a graph $G = G_m$ and M_2 is the partition matroid associated with the colours. Thus for a set of edges E , $r_1(E) = n - K(S)$ where $K(S)$ is the number of components of the graph $G_S = ([w], S)$ and $r_2(S)$ is the number of distinct colours occurring in S . If $i \in W$ then C_i denotes the set of edges of colour i and for $I \subseteq W$, $C_I = \bigcup_{i \in I} C_i$. We will use Edmonds' theorem in the following form:

Theorem 2 *Suppose $|W| = n - 1$. Then a necessary and sufficient condition for the existence of an MST is that*

$$K(d) \leq n - |I| \quad \text{for all } I \subseteq W. \quad (2)$$

[To see this, w.l.o.g. restrict attention in (1) to E_2 of the form C_I and then take $E_1 = W \setminus I$ in (2).]

2 Proof of Theorem 1

Observe first that if $u = u(n) \rightarrow \infty$ slowly, then in a.e. randomly coloured graph process

$$m_C \geq m_0 = \lfloor \frac{1}{2}n(\ln n - \omega) \rfloor \text{ and } m_N \leq m_1 = \lceil n(\ln n + \omega) \rceil.$$

We will start by justifying a concentration on the case $|W| = n - 1$. We will describe a coupled process in which there are never more than $n - 1$ colours used: from $m_{\mathcal{N}}$ onwards, the colours that have not yet been used are randomly changed to one of the $n - 1$ colours that have appeared so far. The relevant properties of this coupled process are

1. For each $m \in [m_0, m_1]$ the edges of G_m are independently randomly coloured from a choice of $n - 1$ colours.
2. If $m_{\mathcal{MT}} > \max\{m_{\mathcal{C}}, m_{\mathcal{N}}\}$ holds for the original process then it also holds for the coupled process.

Thus to prove our theorem we need only prove that

$$\Pr(m_{\mathcal{MT}} > \max\{m_{\mathcal{C}}, m_{\mathcal{N}}\}) = o(1).$$

where \Pr refers to the coupled process.

Fix some m in the range $[m_0, m_1]$. We define the event

$$\mathcal{A}_k = \{\exists I \subseteq W, |I| = k : \kappa(C_I) \geq n - |I| + 1\}.$$

We know that if $|W| = n - 1$, G_m is connected and each colour is used at least once and there is no MST then \mathcal{A}_k occurs for some $k \in [3, n - 2]$ ($\mathcal{A}_1 \cup \mathcal{A}_2$ cannot occur if all $n - 1$ colours are used and \mathcal{A}_{n-1} cannot occur if G_m is connected.) Take a minimal k , corresponding set I and let $S = C_I$.

Claim 1 G_S has no bridges.

Proof If there is a bridge, remove it and all edges of the same colour. Clearly \mathcal{A}_{k-1} occurs, contradicting the minimality of k . \square



With the notation of Claim 1 suppose then that G_S has i isolated vertices and $n - k + x - i$ non-trivial components, $x \geq 1$. Since non-trivial components without bridges have at least three vertices,

$$i + 3(n - k + x - i) \leq n \tag{3}$$

or

$$\begin{aligned} i &\geq n - k + 2x \\ &\geq n - \frac{3}{2}k + \frac{3}{2} \end{aligned}$$

So now let B_k denote the event

$$\{3 \subseteq W, |V| = k, T \subseteq [n] : |T| \leq 3(k-1)/2, \text{ all edges coloured with } / \text{ are contained in } T, \text{ there are } u \geq \max\{fc, i\} \text{ } / \text{-coloured edges}\}.$$

Here T is the set of vertices in the non-trivial components of G_{C_7} . Thus if $|W| = n - 1$,

$$M_m n A_k \subseteq \bigcup_{i=3}^k B_i \quad \text{for } k \geq 3. \tag{4}$$

For $k \geq 9n/10$ we consider a slightly different event.

We first rephrase (2) as

$$K(C_W/J) \leq |V| + 1 \quad \text{for all } J \subseteq W. \tag{5}$$

So if $|W| = n - 1$ and there is no MST then there exist $\ell \geq 1$ colours whose deletion produces $A \geq \ell + 2$ components of sizes n_1, \dots, n_ℓ .

Claim 2 *Some subsequence of the n_i 's sums to between $\ell + 1$ and $n/2$.*

Proof Assume $n \leq \ell \leq \dots \leq A$.

If $n \geq \ell + 1$, one of n_1, \dots, n_{ℓ} and $n_{\ell+1}$ suffices.

Suppose then that $n_i \leq \ell$, $1 \leq i \leq A$.

Choose r such that

$$n_1 \leq n_r \leq n/2, \quad n_{r+1} > n/2$$

and then

$$\begin{aligned} n_{r+1} &> n/2 - n_r \\ &\geq n/2 - \ell \\ &\geq \ell. \end{aligned}$$

and we can take n_1, \dots, n_r .

Note next that if J is minimal in (5) then each colour in J appears at least twice as an edge joining components of $G_{W \setminus J}$.

So if G_m is connected and there is no MST and A_k does not occur for $k \leq 9n/10$ then there is a set L of $1 \leq \ell < n/10$ colours and a set S of size s , $\ell + 1 \leq s \leq n/2$ such that (i) all $t = r(5) = |(S : \bar{S})| \geq 1$ edges are L -coloured, $((5^* : \bar{S})$ is the set of edges joining 5^* and $\bar{S} = V \setminus S$), (ii) the lexicographically first $\max\{2^{\ell-1}, 0\}$ non- $(5' : \bar{S})$ edges joining up components (of the $W \setminus L$ coloured edges) are also L -coloured. Let $T > \ell$ denote this event.

Then

$$\mathcal{C}_m \cap \left(\bigcup_{k=9n/10}^{n-2} \mathcal{A}_k \right) \subseteq \bigcap_{i=1}^{n/10} (J \text{ PT}_m(V_\ell)). \quad (6)$$

It follows from (4) and (6) that

$$\Pr(m_{\mathcal{MT}} > \max\{m_{\mathcal{N}}, m_{\mathcal{C}}\}) \leq o(1) + \sum_{m=m_0}^{m_1} \left[\sum_{k=3}^{9n/10} \Pr_m(\mathcal{B}_k) + \sum_{\ell=2}^{n/10} \Pr_m(\mathcal{D}_\ell) \right] + \Pr \left(\bigcup_{m=m_0}^{m_1} (\mathcal{C}_m \cap \mathcal{A}_{n-2}) \right). \quad (7)$$

Here \Pr_m denotes probability w.r.t. G_m and the $o(1)$ term is the probability that G_{m_0} is connected or that $m_{\mathcal{N}} > m_1$. (Our calculations force us to separate out \mathcal{A}_{n-2} .)

We must now estimate the individual probabilities in (7). It is easier to work with the independent model G_p , $p = m/N$, where each edge occurs independently with probability p and is then randomly coloured. For any event \mathcal{E} we have (see Bollobás [1] Chapter II) the simple bound

$$\Pr_m(\mathcal{E}) \leq 3\sqrt{n \ln n} \Pr_p(\mathcal{E}). \quad (8)$$

where \Pr_p denotes probability w.r.t. the model G_p .

Now, where $p = \alpha \ln n/n$, $1 - o(1) \leq \alpha \leq 2 + o(1)$,

$$\begin{aligned} \Pr_p(\mathcal{B}_k) &\leq \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \binom{n}{t} \binom{n-1}{k} \binom{\binom{t}{2}}{u} \left(1 - \frac{kp}{n-1}\right)^{\binom{n}{2}-u} \left(\frac{kp}{n-1}\right)^u \\ &\leq \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \frac{n^t e^t n^k e^k}{t^t k^k} \left(\frac{t^2 e}{2u}\right)^u n^{-k\alpha(\frac{1}{2}-o(1))} \left(\frac{\alpha k \ln n}{n^2}\right)^u. \end{aligned} \quad (9)$$

Case 1: $3 \leq k \leq k_0 = n/(3 \ln n)$.

$$\begin{aligned} \Pr_p(\mathcal{B}_k) &\leq \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \left(\frac{e^3 n^{1-\alpha(\frac{1}{2}-o(1))}}{k}\right)^k \left(\frac{t}{n}\right)^{2u-t} \left(\frac{\alpha e k \ln n}{2u}\right)^u \\ &= \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \left(\frac{e^3 n^{1-\alpha(\frac{1}{2}-o(1))}}{k}\right)^k \left(\frac{t}{n}\right)^{u-t} \left(\frac{\alpha e k t \ln n}{2un}\right)^u \end{aligned}$$

$$\begin{aligned}
&\leq \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \left(\frac{e^3 n^{1-\alpha(\frac{1}{2}-o(1))} \alpha \epsilon k \ln n}{2kn} \right)^k \left(\frac{t}{n} \right)^{u-t} \left(\frac{\alpha \epsilon k \ln n}{2n} \right)^{u-k} \\
&= O\left(\left(\frac{\ln n}{n^{\frac{1}{2}-o(1)}} \right)^k \right).
\end{aligned}$$

It follows from this and (8) that

$$\begin{aligned}
\sum_{m=m_0}^{m_1} \sum_{k=4}^{k_0} \Pr_m(\mathcal{B}_k) &= O((n \ln n)(\sqrt{n \ln n})((\ln n)^4/n^{2-o(1)})) \\
&= o(1).
\end{aligned} \tag{10}$$

For $k = 3$ we compute $\Pr_m(\mathcal{B}_3)$ directly, but since now $u = t = k = 3$ is forced,

$$\begin{aligned}
\Pr_m(\mathcal{B}_3) &\leq \binom{n}{3}^2 \left(1 - \frac{3}{n-1}\right)^{m-3} \left(\frac{3}{n-1}\right)^3 \frac{\binom{N-3}{m-3}}{\binom{N}{m}} \\
&= O(e^{3\omega} (\ln n)^3 n^{-3/2})
\end{aligned}$$

and so

$$\sum_{m=m_0}^{m_1} \Pr_m(\mathcal{B}_3) = o(1). \tag{11}$$

Case 2: $k_0 < k \leq n/2$.

We now write (9) as

$$\begin{aligned}
\Pr_p(\mathcal{B}_k) &\leq \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \left(\frac{e^3 n^{1-\alpha(\frac{1}{2}-o(1))}}{k} \right)^k \left(\frac{t}{n} \right)^{u-t} \left(\frac{\alpha \epsilon k t \ln n}{2un} \right)^u \\
&\leq \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \left(\frac{e^3 n^{1-\alpha(\frac{1}{2}-o(1))}}{k} \right)^k \left(\frac{t}{n} \right)^{u-t} n^{\frac{\alpha t k}{2n}}
\end{aligned}$$

(after maximising the last term over u)

$$= \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \left(\frac{e^3 n^{1-\frac{\alpha}{2}(1-\frac{t}{n}-\alpha(1))}}{k} \right)^k \left(\frac{t}{n} \right)^{u-t} \quad (12)$$

$$\leq \sum_{t=1}^{3(k-1)/2} \sum_{u=\max\{t,k\}}^{\binom{t}{2}} \left(\frac{e^3 n^{1-\alpha(\frac{1}{8}-\alpha(1))}}{k} \right)^k \quad (13)$$

since $t \leq 3(k-1)/2 \leq 3n/4$.

(13) and (8) clearly imply

$$\sum_{m=m_0}^{m_i} \sum_{k=k_0}^{n/2} \Pr_m(\mathcal{B}_k) = o(1). \quad (14)$$

Case 3: $n/2 < A \leq 9n/10$

Claim 3 Choose any constant $A > 0$. Then, in a.e. process, simultaneously for each $m \in [m_0, m_i]$, the sets of $s \leq A$ vertices of G_m which span at least s edges together contain at most $(\ln n)^{A+1}$ vertices.

Proof We need only prove this for G_{m_i} and since the property is monotone decreasing we need only prove it for G_{p_i} , $p_i = m_i/n$ ([1], Chapter II).
But

$$\begin{aligned} E_{p_i}(\text{number of vertices}) &\leq \sum_{k=3}^A \binom{U}{k} \binom{\binom{k}{2}}{\lambda} P_i^k \\ &= O(e^{2A} (\ln n)^A). \end{aligned}$$

Now use the Markov inequality. Q

It follows that we may rewrite (3) as

$$i + 3(\ln n)^{A+1} + (A+1)(n-k+x-i) \leq n$$

and so

$$\begin{aligned} i &\geq n - \frac{A+1}{J} k - O((\ln n)^{A+1}) \\ &\geq n - \frac{A}{A-1} k. \end{aligned}$$

By making A sufficiently large we see that if $k \leq 9n/10$ then $t \leq 19n/20$ in (12) and consequently

$$\sum_{m=m_0}^{\min\{9n/10, t\}} \sum_{fc=n/2}^{\ell} \hat{m}(B_k) = o(i). \quad (15)$$

Case 4: $A; \geq 9n/10$

$\Pr_p(\mathcal{D}_\ell) \leq$

$$\sum_{s=0}^{n/2} \binom{n}{s} \binom{n-1}{n-s} \frac{s^{(n-s)}}{s!(n-s)!} \left(e_p \mathbf{V} \right)^{\ell - \max\{2\ell-t, 0\}}$$

Let $u(s, t, \ell)$ denote the summand in the above and let $p = a \ln n/n$ and note that $a \in [1 - UJ/\ln n, 2 + u/\ln n]$.

Case 4.1: $i \leq 2\ell$

It will generally be convenient to split s into two ranges:

Case 4.1.1: $s \leq n^{1/10}$

$$\begin{aligned} u(s, \ell, t) &= \binom{n}{s} \binom{n-1}{\ell} \binom{s(n-s)}{t} p^t (1-p)^{s(n-s)-t} \left(\frac{\ell}{n-1} \right)^{2\ell} \\ &\leq \left(\frac{ne}{s} \right)^s \left(\frac{(n-1)e}{\ell} \right)^\ell \left(\frac{s(n-s)e^{1+p\alpha \ln n}}{tn} \right)^t n^{-\alpha s(n-s)/n} \left(\frac{\ell}{n-1} \right)^{2\ell} \\ &\leq \left(\frac{n^{1-\alpha+\alpha s/n} e}{s} \right)^s \left(\frac{\ell e}{n-1} \right)^\ell \left(\frac{e^2 s(n-s) \ln n}{Tn} \right)^t \\ &\leq \left(\frac{n^{1-\alpha+\alpha s/n} e}{s} \right)^s \left(\frac{e^4 s^2 (n-s)^2 (\ln n)^2}{n^3 \ell} \right)^\ell. \end{aligned} \quad (16)$$

Now

$$n^{l-a+as/n} \leq \frac{1}{n^{\alpha}} o(1) e^{\omega} \quad (17)$$

where $a > 1 - u/\ln n$ and $u \rightarrow \infty$ slowly.

So if $s \leq 3e^n$ then (16) implies that

$$u(s, \ell, t) \leq n^{-(1-o(1))\ell},$$

and if $s > 3e^n$

$$\begin{aligned} U(S, J, A) &\leq \left(e^{l - \frac{(s - \ell)^2 (\ln n)^2}{nH}} \right)^\ell \\ &= O\left(\left(\frac{s}{n^{1-o(1)}} \right)^\ell \right). \end{aligned}$$

Case 4.1.2: $s > n^{1/10}$.

Claim 4.1.2.1. For n a.e. process, every G_m , $m \in [m_0, m_j]$ is 5 -w.c. that $TJ(S) \geq \gamma |S| \ln n$ for all $n^{1/10} \leq |S| \leq n/2$; where $\gamma > 0$ is some absolute constant.

Proof (outline) For $|S| \geq n^{2/3}$ one can use the Chernoff bounds on the tails of the binomial $r_j(S)$. If $|S| \leq n^{2/3}$ we use the fact that with high probability (i) G_{m_0} has n^{ϵ} vertices of degree $\leq 6 \ln n$ where $\epsilon = \epsilon(n) \rightarrow 0$ with n , and (ii) in G_{m_i} no set S of size $\leq n/(\ln n)^2$ contains 3151 edges.

So if $s \geq n^{1/10}$ then we can take $t \geq 7s \ln n > 2\ell$ for some constant $7 > 0$ and this case is vacuous.

Case 4.2 : $t > 2\ell$.

$$u(s, \ell, t) \leq \left(\frac{ne}{s} \right)^s \left(\frac{(n-1)e}{\ell} \right)^\ell \left(\frac{s(n-s)e^{1+p\alpha\ell \ln n}}{tn(n-1)} \right)^t n^{-\alpha s(n-s)/n}$$

$$= \left(\frac{n^{1-\alpha+\alpha s/n} e^{\omega}}{s} \right)^s \left(\frac{f(n-l)e^{\omega}}{tn(n-l)} \right)^t \quad (18)$$

Case 4.2.1: $t \leq 2n$ and so $((n-l)e/t)^t \leq (2ne/\ll)^{t/2}$.

$$u(s, \ell, t) \leq \left(\frac{n^{1-\alpha+\alpha s/n} e^{\omega}}{s} \right)^s \left(\frac{20s\ell \ln n}{t^{3/2} n^{1/2}} \right)^t. \quad (19)$$

Case 4.2.1.1: $s < n^{1/10}$. Now (17) gives

$$\begin{aligned} \left(\frac{n^{1-\alpha+\alpha s/n} e^{\omega}}{s} \right)^s &\leq \left(\frac{(1+o(1))e^{\omega+1}}{s} \right)^s \\ &\leq e^{(1+o(1))e^{\omega}} \\ &= e^{\hat{\omega}}, \text{ say,} \end{aligned}$$

and so (19) implies

$$u(s, \ell, t) \leq \left(\frac{e^{\hat{\omega}}}{s^{1-\alpha(1)}} \right)^t. \quad (20)$$

Case 4.2.1.2: $s \geq n^{1/10}$.

Using Claim 4 and (19),

$$u(s, \ell, t) \leq n^{-s/11} \left(\frac{1}{n^{1/2-\alpha(1)} \sqrt{s}} \right)^t.$$

Case 4.2.2: $t \geq 2n$ and so $(ne/\ell)^e \leq e^n \leq e^{t/2}$.

From (18),

$$u(s, \ell, t) \leq \left(\frac{(1+o(1))e^{\omega+1}}{s} \right)^s \left(\frac{20s\ell \ln n}{tn} \right)^t.$$

Case 4.2.2.1: $s < n^{1/10}$.

Arguing as in (20),

$$u(s, \ell, t) \leq \left(\frac{s}{n^{1-o(1)}} \right)^t.$$

Case 4.2.2.2: $s \geq n^{1/10}$.

From Claim 4

$$u(s, \ell, t) \leq \left(\frac{(1+o(1))e^{\omega+1}}{s} \right)^s \left(\frac{A\ell}{n} \right)^t.$$

for some constant $A > 0$. Now this clearly implies

$$u(s, \ell, t) = O(2^{-n}) \tag{21}$$

for $\ell \leq n/(3A)$. For $\ell > n/(3A)$ we have $s \geq \ell$ and

$$u(s, \ell, t) \leq n^{-s/2} A^n$$

and so (21) holds here also.

Summarising,

$$\begin{aligned} \Pr(\mathcal{D}_\ell) &= O \left(\sum_{t=1}^{2\ell} \sum_{s=\ell+1}^{n^{1/10}} \left(\frac{s}{n^{1-o(1)}} \right)^\ell + \sum_{t=2\ell+1}^{2n} \sum_{s=\ell+1}^{n^{1/10}} \left(\frac{s}{n^{\frac{1}{2}-o(1)}} \right)^t \right. \\ &\quad + \sum_{t=2\ell+1}^{2n} \sum_{s=n^{1/10}}^{n/2} \left(\frac{s}{n^{\frac{1}{2}-o(1)}\sqrt{s}} \right)^t + \sum_{s=1}^{n^{1/10}} \sum_{t=2n+1}^{s(n-s)} \left(\frac{s}{n^{\frac{1}{2}-o(1)}} \right)^t \\ &\quad \left. + \sum_{s=n^{1/10}}^{n/2} \sum_{t=2n+1}^{s(n-s)} 2^{-n} \right) \\ &= O(\ell n^{-(.9-o(1))\ell}). \end{aligned}$$

where the double summations correspond to the five cases enumerated above.

Thus, we see that

$$\begin{aligned} \sum_{m=m_0}^{m_1} \sum_{\ell=2}^{n/10} \Pr_m(\mathcal{D}_\ell) &= O((n \ln n)(\sqrt{n \ln n})n^{-1.7}) \\ &= o(1). \end{aligned} \tag{22}$$

We are thus left with $\Pr\left(\bigcup_{m=m_0}^{m_1} (C_m \cap \mathcal{A}_{n-2})\right)$.

We consider G_{m_0} . We know that a.e. G_{m_0} consists of a giant connected component C plus $O(e^\omega)$ isolated vertices T . If $\bigcup_{m=m_0}^{m_1} (C_m \cap \mathcal{A}_{n-2})$ occurs at some time during the process then either

(i) there exist $u, v \in T$ such that the first edges of the process that are incident with each of u and v are the same colour,

OR

(ii) there exists a colour c and a set S , $2 \leq |S| \leq n/2$ such that in G_{m_0} the $t \geq 2$ ($S : \bar{S}$) edges are all of colour c .

(Suppose that deleting the edges of colour c from G_m produces at least three components. If colour k has not occurred by time m_0 then two of these components must be vertices from T , contradicting (i). If G_{m_0} has edges of colour c then deleting these edges must break C into at least three pieces.)

Clearly

$$\Pr((i)) = o(1) + O(e^{2\omega}/n) = o(1).$$

Furthermore

$$\begin{aligned} \Pr_p((ii)) &\leq \sum_{s=2}^{n/2} \binom{n}{s} n^s \sum_{t=2}^{s(n-s)} \binom{s(n-s)}{t} \left(\frac{p}{n}\right)^t (1-p)^{s(n-s)-t} \\ &\leq 2 \sum_{s=2}^{n/2} \binom{n}{s} n^s \sum_{t=2}^{10 \ln n} \frac{(s(n-s))^t}{t!} \left(\frac{\alpha \ln n}{n^2}\right)^t n^{-\alpha s} \\ &\leq n \sum_{s=2}^{n/2} \left(\frac{n^{1-\alpha}}{s}\right)^s \sum_{t=2}^{10 \ln n} \left(\frac{s\alpha \ln n}{n}\right)^t \\ &= O(n^{-(1-o(1))}). \end{aligned}$$



The upper bound is good enough to apply (8) and so $\Pr_{m_0}((ii)) = o(1)$. Thus

$$\Pr \left(\bigcup_{m=m_0}^{m_1} (\mathcal{C}_m \cap \mathcal{A}_{n-2}) \right) = o(1). \quad (23)$$

Our theorem now follows from (7),(10),(11),(14),(15),(22) and (23).

References

- [1] B.Bollobás, *Random graphs*, Academic press, 1985.
- [2] J.Edmonds, *Submodular functions, matroids and certain polyhedra*, in *Combinatorial Structures and their Applications*, R.Guy et al, eds., Gordon and Breach, 1970, pp69-87.

